- Key idea for relaxation techniques intuitive.
- Associate a single equation, corresponding single unknown,  $u_{i,j}$ , with each mesh point in  $\Omega^h$ .
- Then repeatedly "sweep" through mesh, visiting each mesh point in some prescribed order.
- Each time point is visited, adjust value of unknown at grid point so corresponding equation is ("instantaneously") satisfied.
- Adopt a "residual based" approach to locally satisfying the discrete equations.
- Proceed in a manner that generalizes immediately to the solution of *non-linear* elliptic PDEs
- Ignore the treatment of boundary conditions (if conditions are *differential*, will also need to be relaxed)

- Again, adopt "residual-based" approach to the problem of locally satisfying equations via relaxation
- Consider general form of discretized BVP

$$L^h u^h = f^h \tag{1}$$

and recast in canonical form

$$F^h\left[u^h\right] = 0. \tag{2}$$

- Quantity u<sup>h</sup> which appears above is the exact solution of the difference equations.
- Can generally only compute  $u^h$  in the limit of infinite iteration.
- Thus introduce  $\tilde{u}^h$ : "current" or "working" approximation to  $u^h$ , labelling the iteration number by n, and assuming iterative technique *does* converges, have

$$\lim_{n \to \infty} \tilde{u}^h = u^h \tag{3}$$

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• Associated with  $\tilde{u}^h$  is residual,  $\tilde{r}^h$ 

$$\tilde{r}^h \equiv L^h \tilde{u}^h - f^h \tag{4}$$

or in terms of canonical form (2),

$$\tilde{r}^h \equiv F^h \left[ \tilde{u}^h \right] \,. \tag{5}$$

• For specific *component* (grid value) of residual,  $\tilde{r}_{i,j}^h$ , drop the h superscript

$$\tilde{r}_{i,j} = \left[ L^h \tilde{u}^h - f^h \right]_{i,j} \equiv \left[ F^h \left[ u^h \right] \right]_{i,j} \tag{6}$$

• For model problem have

$$\tilde{r}_{i,j} = h^{-2} \left( \tilde{u}_{i+1,j} + \tilde{u}_{i-1,j} + \tilde{u}_{i,j+1} + \tilde{u}_{i,j-1} - 4\tilde{u}_{i,j} \right) - f_{i,j}$$
(7)

• Relaxation: adjust  $\tilde{u}_{i,j}$  so corresponding residual is "instantaneously" zeroed

- Useful to appeal to Newton's method for single non-linear equation in a single unknown.
- In current case, difference equation is *linear* in  $\tilde{u}_{i,j}$ : can *solve* equation with single Newton step.
- However, can also apply relaxation to *non-linear* difference equations, then can only zero residual in limit of infinite number of Newton steps.
- Thus write relaxation in terms of update,  $\delta \tilde{u}_{i,j}^{(n)}$ , of unknown

$$\tilde{u}_{i,j}^{(n)} \longrightarrow \tilde{u}_{i,j}^{(n+1)} = \tilde{u}_{i,j}^{(n)} + \delta \tilde{u}_{i,j}^{(n)}$$
(8)

where, again, (n) labels the iteration number.

• Using Newton's method, have

$$\widetilde{u}_{i,j}^{(n+1)} = \widetilde{u}_{i,j}^{(n)} - \widetilde{r}_{i,j} \left[ \frac{\partial F_{i,j}^{h}}{\partial u_{i,j}} \Big|_{u_{i,j} = \widetilde{u}_{i,j}^{(n)}} \right]^{-1}$$

$$= \widetilde{u}_{i,j}^{(n)} - \frac{\widetilde{r}_{i,j}}{-4h^{-2}}$$

$$= \widetilde{u}_{i,j}^{(n)} + \frac{1}{4}h^{2}\widetilde{r}_{i,j}$$
(10)
(11)

- Precise computation of the residual needs clarification.
- Iterative method takes an entire vector of unknowns  $\mathbf{u}^{(n)}$  to new estimate  $\mathbf{u}^{(n+1)}$ , but works on a component by component basis.

- In computing individual residuals, could either choose only "old" values; i.e. values from iteration n, or, wherever available, could use "new" values from iteration n + 1, with the rest from iteration n.
- First approach is known as *Jacobi relaxation*, residual computed as

$$\tilde{r}_{i,j} = h^{-2} \left( \tilde{u}_{i+1,j}^{(n)} + \tilde{u}_{i-1,j}^{(n)} + \tilde{u}_{i,j+1}^{(n)} + \tilde{u}_{i,j-1}^{(n)} - 4\tilde{u}_{i,j}^{(n)} \right) - f_{i,j}$$
(12)

Second approach is known as Gauss-Seidel relaxation: assuming lexicographic ordering of unknowns, i = 1, 2, ... n, j = 1, 2, ... n, i index varies most rapidly, residual is

$$\tilde{r}_{i,j} = h^{-2} \left( \tilde{u}_{i+1,j}^{(n)} + \tilde{u}_{i-1,j}^{(n+1)} + \tilde{u}_{i,j+1}^{(n)} + \tilde{u}_{i,j-1}^{(n+1)} - 4\tilde{u}_{i,j}^{(n)} \right) - f_{i,j}$$
(13)

• Make few observations/comments about these two relaxation methods.

- At each iteration "visit" each/every unknown *exactly once*, modifying its value so that local equation is instantaneously satisfied.
- Complete pass through the mesh of unknowns (i.e. a complete iteration) is known as a *relaxation sweep*.
- For Jacobi, visit order clearly irrelevant to what values are obtained at end of each iteration
- Fact is advantageous for parallelization but storage is required for *both* the new and old values of  $\tilde{u}_{i,j}$ .
- For Gauss-Seidel (GS), need storage for current estimate of  $\tilde{u}_{i,j}$ : sweeping order *does* impact the details of the solution process.
- Thus, lexicographic ordering does not parallelize, but for difference equations such as those for model problem, with nearest-neighbour couplings, so called *red-black* ordering *can* be parallelized, has other advantages.
- Red-black ordering: appeal to the red/black squares on a chess board—two subiterations
  - 1. Visit and modify all "red" points, i.e. all (i, j) such that mod(i + 1, 2) = 0
  - 2. Visit and modify all "black" points, i.e. all (i, j) such that mod(i + 1, 2) = 1

- Key issue with any iterative technique is whether or not the iteration will converge
- For Gauss-Seidel and related, was examined comprehensively in the 1960's and 1970's: won't go into much detail here
- Very rough rule-of-thumb: GS will converge if linear system *diagonally dominant*
- Write (linear) difference equations

$$L^h u^h = f^h \tag{14}$$

in matrix form

$$\mathbf{A}\mathbf{u} = \mathbf{b} \tag{15}$$

where **A** is an  $N \times N$  matrix ( $N \equiv$  total number of unknowns); **u** and **b** are N-component vectors

7 7

• A diagonally dominant if

$$|a_{ij}| \le \sum_{j=1, j \ne i}^{N} |a_{ij}|, i = 1, 2, \cdots N$$
 (16)

with strict inequality holding for at least one value of i.

- Another practical issue: *how* to monitor convergence?
- For relaxation methods, two natural quantities to look at: the *residual norm*  $\|\tilde{r}^h\|$  and the *solution update norm*  $\|\tilde{u}^{(n+1)} \tilde{u}^{(n)}\|$
- Both should tend to 0 in limit of infinite iteration, for convergent method be convergent.
- In practice, monitoring residual norm straightforward and often sufficient.

- Consider the issue of convergence in more detail.
- For a general iterative process leading to solution vector  $\mathbf{u}$ , and starting from some initial estimate  $\mathbf{u}^{(0)}$ , have

$$\mathbf{u}^{(0)} \to \mathbf{u}^{(1)} \to \cdots \to \mathbf{u}^{(n)} \to \mathbf{u}^{(n+1)} \to \cdots \to \mathbf{u}$$
 (17)

• For the residual vector, have

$$\mathbf{r}^{(0)} \to \mathbf{r}^{(1)} \to \cdots \to \mathbf{r}^{(n)} \to \mathbf{r}^{(n+1)} \to \cdots \to \mathbf{0}$$
 (18)

• For the solution error,  $\mathbf{e}^{(n)} \equiv \mathbf{u} - \mathbf{u}^{(n)}$ , have

$$\mathbf{e}^{(0)} \to \mathbf{e}^{(1)} \to \dots \to \mathbf{e}^{(n)} \to \mathbf{e}^{(n+1)} \to \dots \to \mathbf{0}$$
 (19)

- For linear relaxation (and basic idea generalizes to non-linear case), can view transformation of error, e<sup>(n)</sup>, at each iteration in terms of linear matrix (operator), G, known as the *error amplification matrix*.
- Then have

$$\mathbf{e}^{(n+1)} = \mathbf{G}\mathbf{e}^{(n)} = \mathbf{G}^2\mathbf{e}^{(n-1)} = \dots = \mathbf{G}^n\mathbf{e}^{(0)}$$
(20)

• Asymptotically convergence determined by the *spectral radius*,  $\rho(\mathbf{G})$ :

$$\rho(\mathbf{G}) \equiv \max_{i} |\lambda_i(\mathbf{G})| \tag{21}$$

and the  $\lambda_i$  are the eigenvalues of **G**.

• That is, in general have

$$\lim_{n \to \infty} \frac{\|\mathbf{e}^{(n+1)}\|}{\|\mathbf{e}^n\|} = \rho(\mathbf{G})$$
(22)

• Can then define asymptotic convergence rate, R

$$R \equiv \log_{10} \left( \rho^{-1} \right) \tag{23}$$

- Useful interpretation of R:  $R^{-1}$  is number of iterations needed asympotically decrease  $\|\mathbf{e}^{(n+1)}\|$  by an order of magnitude.
- Now consider the computational cost of solving discrete BVPs using relaxation.
- "Cost" is to be roughly identified with CPU time, and want to know how cost scales with key numerical parameters
- Here, assume that there is *one* key parameter: total number of grid points, N
- For concreteness/simplicity assume following
  - 1. The domain  $\Omega$  is d-dimensional (d = 1, 2, 3 being most common)
  - 2. There are n grid points per edge of the (uniform) discrete domain,  $\Omega^h$
- Then total number of grid points, N, is

$$N = n^d \tag{24}$$

 Further define the computational cost/work, W ≡ W(N) to be work required to reduce the error, ||e|| by an order of magnitude; definition suffices to compare methods

- Clearly, best case situation is W(N) = O(N).
- For Gauss-Seidel relaxation, state without proof that for model problem (and for many other elliptic systems in d = 1, 2, 3), have

$$\rho\left(\mathbf{G}_{\mathrm{GS}}\right) = 1 - O(h^2) \tag{25}$$

• Implies that relaxation work,  $W_{\rm GS}$ , required to reduce the solution error by order of magnitude is

$$W_{\rm GS} = O\left(h^{-2}\,{\rm sweeps}\right) = O\left(n^2\,{\rm sweeps}\right)$$
 (26)

• Each relaxation sweep costs  $O(n^d) = O(N)$ , so

$$W_{\rm GS} = O\left(n^2 N\right) = O\left(N^{2/d} N\right) = O\left(N^{(d+2)/d}\right) \tag{27}$$

• Tabulating the values for d = 1, 2 and 3



- Scaling improves as d, increases, but  $O(N^2)$  and  $O(N^{5/3})$  for the cases d = 2 and d = 3 are pretty bad
- Reason that Gauss-Seidel is rarely used in practice, particularly for large problems (large values of *n*, small values of *h*).

# Successive Over Relaxation (SOR)

- Historically, researchers studying Gauss-Seidel found that convergence of the method could often be substantially improved by systematically "over correcting" the solution, relative to what the usual GS computation would give
- Algorithmically, change from GS to SOR is very simple, and is given by

$$\tilde{u}_{i,j}^{(n+1)} = \omega \,\hat{u}_{i,j}^{(n+1)} + (1-\omega) \,\tilde{u}_{i,j}^{(n)} \tag{28}$$

- Here,  $\omega$  is the *overrelaxation parameter*, typically chosen in the range  $1 \leq \omega < 2$ , and  $\hat{u}_{i,j}^{(n+1)}$  is the value computed from the normal Gauss-Seidel iteration.
- When  $\omega = 1$ , recover GS iteration itself: for large enough  $\omega$ , e.g.  $\omega \ge 2$ , iteration will become unstable.
- At best, the use of SOR reduces number of relaxation sweeps required to drop the error norm by an order of magnitude to O(n):

$$W_{\text{SOR}} = O\left(nN\right) = O\left(N^{1/d}N\right) = O\left(N^{(d+1)/d}\right)$$
(29)

# Successive Over Relaxation (SOR)

• Again tabulating the values for d = 1, 2 and 3 find



- Thus note that optimal SOR is not unreasonable for "moderately-sized" d = 3 problems.
- Key issue: How to choose  $\omega = \omega_{\rm OPT}$  in order to get optimal performance
- Except for cases where  $\rho_{GS} \equiv \rho(G_{GS})$  is known explicitly,  $\omega_{OPT}$  needs to be determined *empirically* on a case-by-case *and* resolution-by-resolution basis.
- If  $ho_{\mathrm{GS}}$  is known, then one has

$$\omega_{\rm OPT} = \frac{2}{1 + \sqrt{1 - \rho_{\rm GS}}} \tag{30}$$

- Slow convergence rate of relaxation methods such as Gauss-Seidel → not much use for solving discretized elliptic problems
- Now consider an even simper model problem to show what relaxation *does* tend to do well
- Will elucidate why relaxation is so essential to the multi-grid method.
- Consider a one dimensional (d = 1) "elliptic" model problem, i.e. an *ordinary* differential equation to be solved as a two-point boundary value problem.
- Equation to be solved is

$$Lu(x) \equiv \frac{d^2u}{d^2x} = f(x) \tag{31}$$

where f(x) is specified source function.

• Solve (31) on a domain,  $\Omega$ , given by

$$\Omega \,:\, 0 \le x \le 1 \tag{32}$$

subject to the (homogeneous) boundary conditions

$$u(0) = u(1) = 0 \tag{33}$$

• To discretize this problem, introduce a uniform mesh,  $\Omega^h$ 

$$\Omega^{h} = \{ (i-1)h, \ i = 1, 2, \cdots n \}$$
(34)

where  $n = h^{-1} + 1$  and h, as usual, is the mesh spacing.

• Now discretize (31) to  $O(h^2)$  using usual FDA for second derivative, getting

$$\frac{u_{i+1} - 2u_i + u_{i-1}}{h^2} = f_i, \ i = 2, 3, \dots n - 1$$
(35)

• Last equations, combined with boundary conditions

$$u_1 = u_n = 0 \tag{36}$$

yield set of n linear equations in n unknowns,  $u_i, i = 1, 2, \dots n$ .

• For analysis, useful to eliminate  $u_1$  and  $u_n$  from the discretization , so that in terms of generic form

$$L^h u^h = f^h \tag{37}$$

 $u^h$  and  $f^h$  are n-2-component vectors;  $L^h$  is an  $(n-2) \times (n-2)$  matrix.

- Consider a specific relaxation method, known as *damped Jacobi*, chosen for ease of analysis, but representative of other relaxation techniques, including Gauss-Seidel.
- Using vector notation, damped Jacobi iteration is given by

$$\tilde{\mathbf{u}}^{(n+1)} = \tilde{\mathbf{u}}^{(n)} - \omega \mathbf{D}^{-1} \tilde{\mathbf{r}}^{(n)}$$
(38)

where  $\tilde{\mathbf{u}}$  is the approximate solution of the difference equations,  $\tilde{\mathbf{r}}$  is the residual vector,  $\omega$  is an (underrelaxation) parameter, and  $\mathbf{D}$  is the main diagonal of  $L^h_{\frac{18}{18}}$ .

• For model problem, have

$$\mathbf{D} = -2h^{-2}\mathbf{1} \tag{39}$$

- Note that when  $\omega = 1$ , this is just the usual Jacobi relaxation method discussed previously.
- Examine the effect of (38) on residual vector,  $\tilde{\mathbf{r}}$ .

Have

$$\tilde{\mathbf{r}}^{(n+1)} \equiv L^h \tilde{\mathbf{u}}^{(n+1)} - f^h$$
 (40)

$$= L^{h} \left( \tilde{\mathbf{u}}^{(n)} - \omega \mathbf{D}^{-1} \tilde{\mathbf{r}}^{(n)} \right) - f^{h}$$
(41)

$$= (\mathbf{1} - \omega L^h \mathbf{D}^{-1}) \tilde{\mathbf{r}}^{(n)}$$
(42)

$$\equiv \mathbf{G}_{\mathrm{R}}\tilde{\mathbf{r}}^{(n)} \tag{43}$$

where the residual amplification matrix,  $\mathbf{G}_{\mathrm{R}},$  is

$$\mathbf{G}_{\mathrm{R}} \equiv \mathbf{1} - \omega L^h \mathbf{D}^{-1} \tag{44}$$

• Have following relation for residual at the *n*-th iteration,  $\tilde{\mathbf{r}}^{(n)}$ , in terms of initial residual,  $\tilde{\mathbf{r}}^{(0)}$ :

$$\tilde{\mathbf{r}}^{(n)} = \left(\mathbf{G}_{\mathrm{R}}\right)^{n} \, \tilde{\mathbf{r}}^{(0)} \tag{45}$$

- Now,  $G_R$ , has complete set of orthogonal eigenvectors,  $\phi_m$ ,  $m = 1, 2, \dots n 2$  with corresponding eigenvalues  $\mu_m$ .
- Thus, can expand initial residual,  $\tilde{\mathbf{r}}^{(0)}$ , in terms of the  $\phi_m$ :

$$\tilde{\mathbf{r}}^{(0)} = \sum_{m=1}^{n-2} c_m \phi_m \tag{46}$$

where the  $c_m$  are coefficients.

• Immediately yields following expansion for the residual at the *n*-th iteration:

$$\tilde{\mathbf{r}}^{(n)} = \sum_{m=1}^{n-2} c_m \left(\mu_m\right)^n \phi_m.$$
(47)

- Thus, after n sweeps, the m-th "Fourier" component of the initial residual vector,  $\tilde{r}^{(0)}$  is reduced by a factor of  $(\mu_m)^n$ .
- Now consider the specific value of the underrelaxation parameter,  $\omega = 1/2$ .
- Left as exercise to verify that eigenvectors,  $\phi_m$  are

 $\phi_m = [\sin(\pi mh), \sin(2\pi mh), \cdots, \sin((n-2)\pi mh)]^T, \ m = 1, 2, \cdots, n-2$ (48)

while eigenvalues,  $\mu_m$ , are

$$\mu_m = \cos^2\left(\frac{1}{2}\pi mh\right), \ m = 1, 2, \cdots, n-2$$
(49)

- Note that each eigenvector,  $\phi_m$ , has associated ''wavelength'',  $\lambda_m$ 

$$\sin\left(\pi mx\right) = \sin\left(2\pi x/\lambda_m\right) \,. \tag{50}$$

Thus

$$\lambda_m = \frac{2}{m} \tag{51}$$

- As m increases (and  $\lambda_m$  decreases), "frequency" of  $\phi_m$  increases, and, from (49), the eigenvalues,  $\mu_m$ , decrease in magnitude.
- Asymptotic convergence rate of the relaxation scheme is determined by largest of the  $\mu_m$ ,  $\mu_1$

$$\mu_1 = \cos^2\left(\frac{1}{2}\pi h\right) = 1 - \frac{1}{4}\pi^2 h^2 + \dots = 1 - O(h^2).$$
(52)

- Implies that  $O(n^2)$  sweeps are needed to reduce norm of residual by order of magnitude.
- Thus see that asymptotic convergence rate of relaxation scheme is dominated by the (slow) convergence of *smooth* (low frequency, long wavelength) components of the residual,  $\tilde{\mathbf{r}}^{(n)}$ , or, equivalently, the solution error,  $\tilde{\mathbf{e}}^{(n)}$ .
- Comment applies to many other relaxation schemes applied to typical elliptic problems, including Gauss-Seidel.

# **Definition of "High Frequency"**



- Illustration of the definition of "high frequency" components of the residual or solution error, on a fine grid,  $\Omega^h$ , relative to a coarse grid,  $\Omega^{2h}$ .
- As illustrated in the figure, any wave with λ < 4h cannot be represented on the coarse grid (i.e. the Nyquist limit of the coarse grid is 4h.)

- Highly instructive to consider what happens to high frequency components of the residual (or solution error) as the damped Jacobi iteration is applied.
- Per previous figure, are concerned with eigenvectors,  $\phi_m$ , such that

$$\lambda_m \le 4h \longrightarrow mh \ge \frac{1}{2} \tag{53}$$

• In this case, have

$$\mu_m = \cos^2\left(\frac{1}{2}\pi mh\right) \le \cos^2\left(\frac{\pi}{4}\right) \le \frac{1}{2} \tag{54}$$

- Thus, all components of residual (or solution error) that cannot be represented on  $\Omega^{2h}$  get supressed by a factor of at least 1/2 per sweep.
- Furthermore, rate at which high-frequency components are liquidated is *independent* of the mesh spacing, *h*.

#### • SUMMARY:

- Relaxation tends to be a *dismal solver* of systems  $L^h u^h = f^h$ , arising from FDAs of elliptic PDEs.
- But, tends to be a *very good smoother* of such systems—crucial for the success of multi-grid method.

• Again consider model problem

$$u(x,y)_{xx} + u_{yy} = f(x,y)$$
(55)

solved on the unit square subject to homogeneous Dirichlet BCs

- Again discretize system using standard 5-point  $O(h^2)$  FDA on a uniform  $n\times n$  grid with mesh spacing h
- Yields algebraic system  $L^h u^h = f^h$ .
- Assume system solved with iterative solution technique: start with some initial estimate  $\tilde{\mathbf{u}}^{(0)}$ , iterate until  $\|\tilde{\mathbf{r}}^{(n)}\| \leq \epsilon$ .
- Several questions concerning solution of the discrete problem then arise:
  - 1. How do we choose n (equivalently, h)?
  - 2. How do we choose  $\epsilon$ ?
  - 3. How do we choose  $\tilde{\mathbf{u}}^{(0)}$ ?
  - 4. How fast can we "solve"  $L^h u^h = f^h$ ?
- Now provide at least partial answers to all of these questions.

#### 1) Choosing n (equivalently h)

- Ideally, would choose h so that error  $\|\tilde{\mathbf{u}} \mathbf{u}\|$ , satisfies  $\|\tilde{\mathbf{u}} \mathbf{u}\| < \epsilon_u$ , where  $\epsilon_u$  is a user-prescribed error tolerance.
- From discussions of FDAs for time-dependent problems, already know how to estimate solution error for essentially *any* finite difference solution.
- Assume that FDA is centred and  $O(h^2)$
- Richardson tells us that for solutions computed on grids with mesh spacings h and 2h, expect

$$u^h \sim u + h^2 e_2 + h^4 e_4 + \cdots$$
 (56)

$$u^{2h} \sim u + 4h^2 e_2 + 16h^4 e_4 + \cdots$$
 (57)

• So to leading order in h, error,  $u^h - u$ , is

$$e \sim \frac{u^{2h} - u^h}{3h^2} \tag{58}$$

• Thus, basic strategy is to perform *convergence tests*—comparing finite difference solutions at different resolutions, increase (or decrease) h until the level of error is satisfactory.

## 2) Choosing $\epsilon$ ( $\epsilon_r$ )

• Consider following 3 expressions:

$$L^h u^h - f^h = 0 (59)$$

$$L^h \tilde{u}^h - f^h = \tilde{r}^h \tag{60}$$

$$L^h u - f^h = \tau^h \tag{61}$$

where  $u^h$  is exact solution of finite difference equations,  $\tilde{u}^h$  is approximate solution of the FDA, u is the (exact) solution of the continuum problem, Lu - f = 0.

 (59) is simply our FDA written in a canonical form, (60) defines the residual, while (61) defines the truncation error.

 Comparing (60) and (61), see that it is natural to stop iterative solution process when

$$\|\tilde{r}^h\| \sim \|\tau^h\| \tag{62}$$

- Leaves problem of estimating size of truncation error
- Will see later how estimates of  $\tau^h$  arise naturally in multi-grid algorithms.

# 3) Choosing $\tilde{\mathbf{u}}^{(0)}$

- Key idea is to use solution of coarse-grid problem as initial estimate for fine-grid problem.
- Assume that have determined satisfactory resolution h; i.e. wish to solve

$$L^h u^h = f^h \tag{63}$$

- Then first pose and solve (perhaps approximately) corresponding problem (i.e. same domain, boundary conditions and source function, f) on mesh with spacing 2h.
- That is, solve

$$L^{2h}u^{2h} = f^{2h} (64)$$

then set initial estimate,  $(u^h)^{(0)}$  via

$$(u^{h})^{(0)} = \bar{I}^{h}_{2h} \, u^{2h} \tag{65}$$

- Here  $\overline{I}_{2h}^h$  is known as a *prolongation operator* and transfers a coarse-grid function to a fine-grid.
- Typically,  $\bar{I}_{2h}^h$  will perform *d*-dimensional polynomial interpolation of the coarse-grid unknown,  $u^{2h}$ , to a suitable order in the mesh spacing, *h*.
- Chief advantage of approach is that solution of coarse-grid problem should be inexpensive to solve relative to fine-grid problem
- Specifically, cost of solving on the coarse-grid should be no more than 2<sup>-d</sup> of cost of soving the fine-grid equations.
- Furthermore, can apply basic approach *recursively*: initialize  $u^{2h}$  with  $u^{4h}$  result, initialize  $u^{4h}$  with  $u^{8h}$  etc
- Thus lead to consider a general *multi-level* technique for treatment of discretized elliptic PDEs
- Solution of a fine-grid problem is preceded by solution of series of coarse-grid problems.

- Label each distinct mesh spacing with an integer  $\ell$ 

$$\ell = 1, 2, \cdots \ell_{\max} \tag{66}$$

where  $\ell = 1$  labels coarsest spacing  $h_1$ ,  $\ell = \ell_{\max}$  labels finest spacing  $h_{\ell_{\max}}$ .

• Almost always most convenient (and usually most efficient) to use mesh spacings,  $h_{\ell}$ , satisfying

$$h_{\ell+1} = \frac{1}{2}h_{\ell} \tag{67}$$

• This implies

$$n_{\ell+1} \sim 2^d n_\ell \,. \tag{68}$$

- Will assume in following that multi-level algorithms use 2:1 mesh hierarchies
- Use ℓ itself to denote resolution associated with a particular finite difference approximation, or transfer operator.
- That is, define  $u^\ell$  via

$$u^{\ell} = u^{h_{\ell}} \tag{69}$$

• Could then rewrite (65) as

$$u^{\ell+1} = \bar{I}_{\ell}^{\ell+1} u^{\ell}$$
 (70)

• Pseudo-code for general multi-level technique:

```
for \ell = 1, \ell_{\max}

if \ell = 1 then

u^{\ell} := 0

else

u^{\ell} := \overline{I}_{\ell-1}^{\ell} u^{\ell-1}

end if

solve_iteratively (u^{\ell})

end for
```

- Very simple and intuitive algorithm: when novices hear term "multi-grid", sometimes think that this sort of approach is all that is involved.
- However, multi-grid algorithm uses grid hierarchy for entirely different, and more essential purpose!

## 4) How fast can we solve $L^h u^h = f^h$ ?

- Answer to this question is short and sweet.
- Properly constructed multi-grid algorithm can solve discretized elliptic equation in  ${\cal O}(N)$  operations!
- This is, of course, computationally optimal: main reason that multi-grid methods have become such an important tool for numerical analysts—numerical relativists included.